

Measuring the DNS

Seminar Internet Measurements, TU-Berlin

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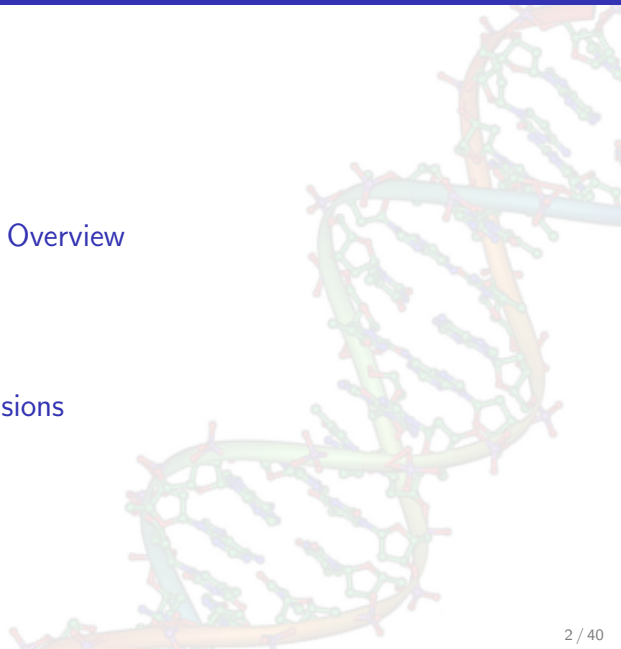
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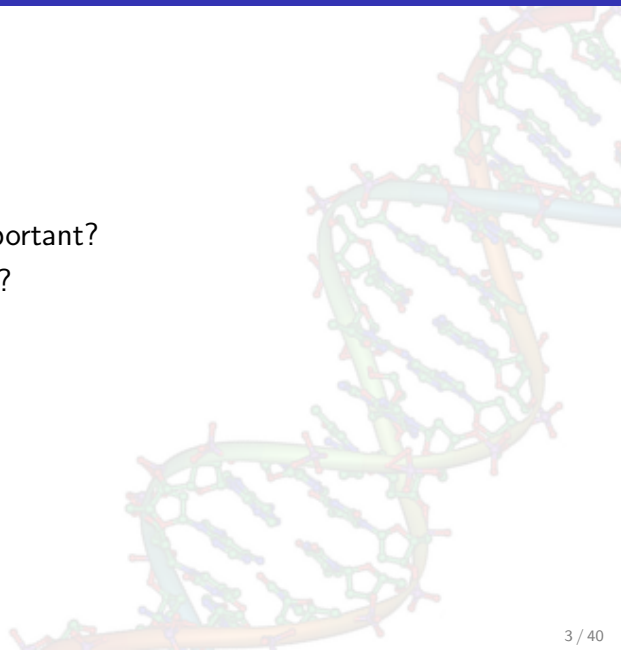
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Motivation

- Why is DNS so important?
- What is it good for?
- Who uses it?



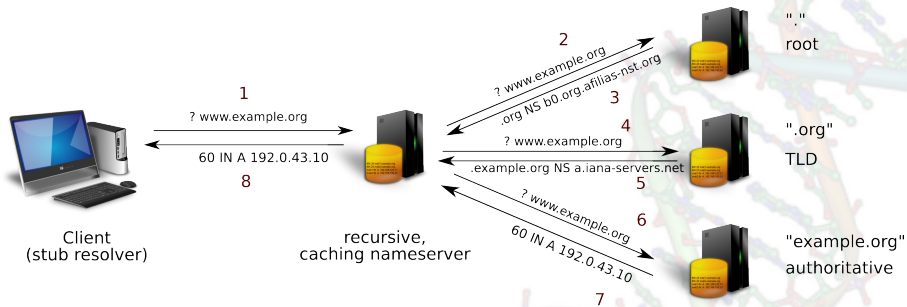
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 - Technical Introduction
 - DNS - mode of operation
 - DNS-(Ab)uses today
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Technically speaking...

- DNS maps hostnames to IP-addresses and vice versa
- until 1983: hosts file
- RFC 882 and RFC 883
- not only address mapping
- data organized in Resource-Records
- different types of RR exist: A, MX, TXT, NS, SOA, ...
- a DNS query consists of a chain of sections
- usually transmitted via UDP
- TCP is used for big queries, e.g. zone transfers
- characterization: globally distributed database

A simple DNS-Lookup

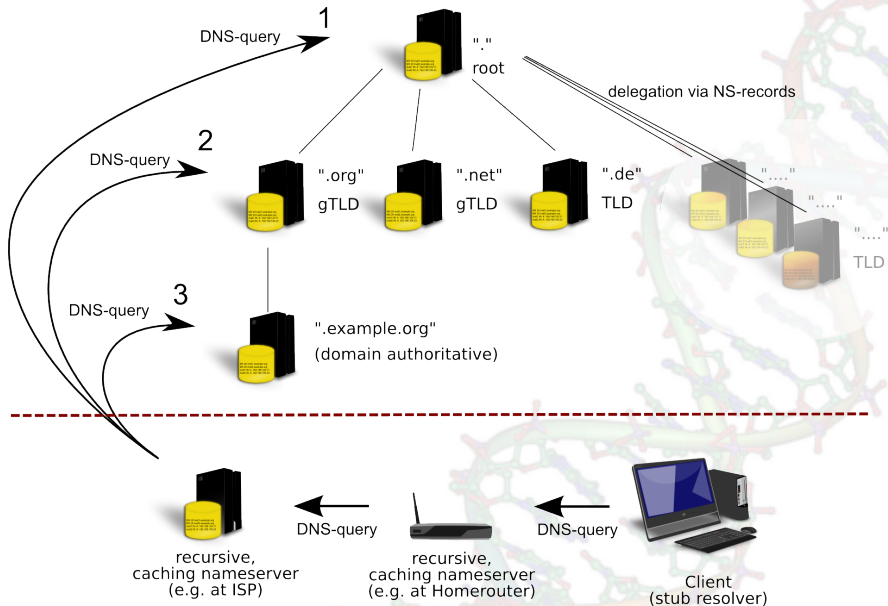


How DNS works

- client: stub resolver
- asks dedicated server
- often embedded in DSL Router
- usually forwards query to ISP
- builds a hierarchy of nameservers
- delegations exist between root, TLDs, domains, ...

(see next slide)

The DNS hierarchy



DNS-(Ab)uses today

- CDNs¹: Records with low TTL
- used for traffic management, multiple levels of indirection
- ssh-key-fingerprints
- google site verification via TXT RRs
- Anti-Spam solutions like SPF, also TXT

⇒ partly considered an abuse, e.g. by P. Vixie [7]

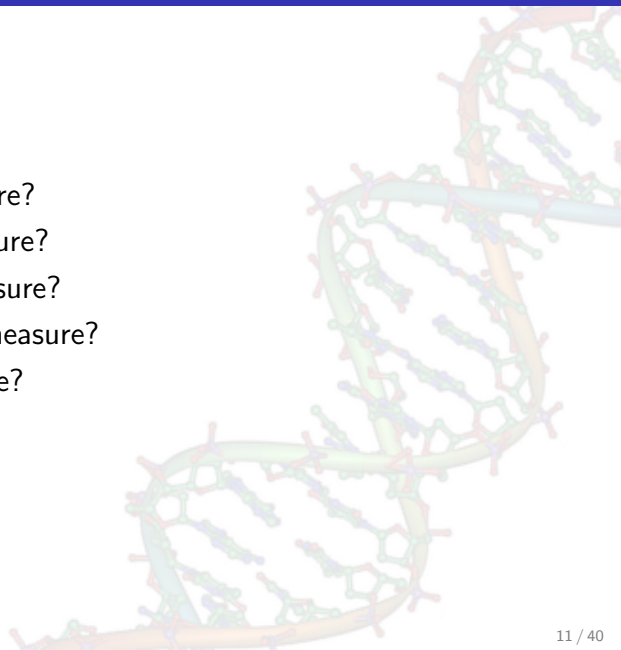
¹Content Distribution Networks

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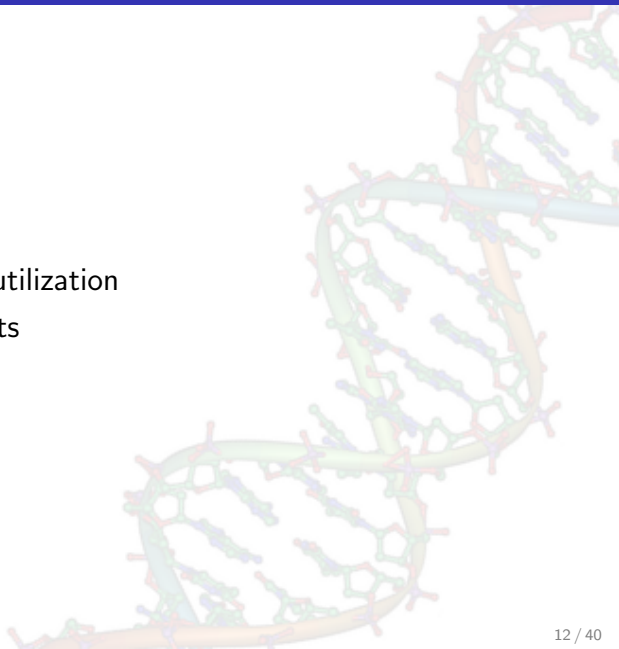
Considerations beforehand

- How can we measure?
- What can we measure?
- Where can we measure?
- Where should we measure?
- Why do we measure?



What to measure

- response time
- packet size
- number of queries
- cache hit rate and utilization
- valid/invalid requests
- success/failure rate
- types of queries
- ...



Where to measure

- application level (e.g. response time)
- client machine (e.g. number of failed/successful queries)
- customer router (e.g. cache utilization, response time, ...)
- provider (e.g. cache hit rate, failed/successful queries,...)
- root zone (e.g. illegal requests, wrong TLDs, ...)
- TLD (e.g. typo domains, number of clients,...)
- 2nd level domain (e.g. dos attacks, v4/v6 ratio, ...)

Levels in the DNS

- stub resolver at the client
- customer cache in the broadband router
- server at the ISP or company
- domain server
- TLD servers
- root servers
- New in the game: dedicated public DNS providers (Google, OpenDNS)

⇒ Each level is special in the data observable!

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Client Level (I)

Considerations:

- at the bottom level
- maximum 'propagation delay':
- requests have to travel through whole DNS
- all layers are caching
- impact of these caches?

Client Level (II)

“DNS Performance and the Effectiveness of Caching” [5]:

- traces collected at two large universities
- 23% of the queries returned no answer
 - responsible for more than half of the DNS-packets:
 - retransmission of failed requests by clients!
- most frequently requested query types:
 - A records (60%), Hostnames to Addresses
 - PTR, reverse lookups (24% to 31%)

Client Level (III)

“DNS Performance and the Effectiveness of Caching” [5]:

- latency of queries
 - number of referrals: big impact
 - for every referral another nameserver has to be contacted
- 70% returned direct answers from cache
- impact of caching can be observed (latency)
- between 10% and 42% result in a negative answer,
- explanation: no reverse-mapping, server misconfiguration, other failures

Client Level (IV)

“DNS Performance and the Effectiveness of Caching” [5]:

- idea: negative caching could have impact on overall performance
- but: heavy-tailed distribution of names
- the effect of negative caching is limited

- idea: negative caching could have impact on overall performance
- again: heavy-tailed nature of access to names!
 - but TTL for NS-records is reasonable high (cached)
 - use of dynamic low-TTL A-record bindings should not degrade DNS performance
- disagrees with Vixie in [7]

“An analysis of wide-area name server traffic[...]” [4]:

Considerations:

- aggregates queries from several users
- currently their location is used for geolocation
- traditionally: fast servers with huge caches

Observations:

- cache hit rate was over estimated (last slides)
- but can serve as 'firewall' between WAN/root DNS and malicious or badly broken programs/users

TLD-Level (I)

- 1 of 13 nameservers for .com and .net (g.gtld-servers.net)
- measurements below the rootservers
- passive measurements (packet dumps)

Expectations by Osterweil et al. [6]:

- audience for the generic TLDs is not country-specific
- resolvers will probe each authoritative server of a zone over time
- nearly all resolvers in the Internet should be seen
- should also see more variance than a rootserver

TLD-Level (II)

Observations:

- traffic from the whole active IP-Range of the Internet is being observed,
- most prominently asked: A (70%)², AAAA (15%) and MX (10%).
- also deprecated types are still being queried, e.g. A6 records, superseded by AAAA

²all approximately values derived from a figure

TLD-Level (III): Top Talkers

Main Aspect of their work:

- small set of 40.000 resolvers, called top-talkers, is responsible for 90% of overall traffic
- the set of top-talkers is highly evolving
- the same set of clients only accounts for 84% percent of the queries after 10 days,
- a huge number of new clients is constantly detected.
- list of top-talkers can be used as a filter in large-scale measurements to dramatically reduce the amount of datapoints, with small effect on accuracy of the measurements.
- yet unclear if a global set of top-talkers can be identified

Informal summary: “Heavy tail is everywhere.”

Root-Level (I)

Considerations:

- up-most level in the hierarchy of nameservers.
- these servers act as last resort if no authority can be found for a subtree
- hit by a lot of attacks (actively and by misconfigurations)
- key element of the Internet (see latest discussions on control over them)
- without them as single root: no anchor, no determinism
- when they fail the Internet will break for end users after caches time out

Root-Level (II)

“DNS measurements at a root server” [2]

- massive amount of erroneous queries
- not always clear, if queries are result of a misconfiguration or attack
- “60-85% of observed queries were repeated from the same host within the measurement interval.”
- “Over 14% of a root server’s query load is due to queries that violate the DNS specification.”
- abused for amplification/reflection attacks using forged source addresses (UDP)

Root-Level (III)

Castro et. al. in [3]

- “an estimated 98% of the traffic at the root servers should not be there at all”
- very high level of DNS pollution.

main findings:

- most of the traffic could be avoided by proper **caching**
- huge amount of traffic originates at **improperly configured** resolvers
- outstanding example:
 - one client querying a non-existent SOA record over 2 million times in an hour
 - caching would have reduced this to less than 10.

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- pseudo-RR was needed, because header had no bits left
- packets getting bigger
- sometimes too big for UDP
- usage of EDNS is increasing:
 - DNSSEC (cryptographically signed DNS data)
 - client IP transmission (for geolocation)

DNSSEC - What and Why

- provides integrity for all replies
- Why? SSH-key fingerprints, e-mail transportation, Online Banking, Domain stealing (ebay)
- each record has to be signed individually, as each one is individually requested.
- cryptographic algorithms are used must be compatible with caching.

DNSSEC - Impacts?

Considerations:

- operation will become more complex, making human error more likely
- bandwidth used will increase because signatures and keys have to be transmitted
- caches will have to hold more information
- validating resolvers will have to do complex cryptographic operations to validate each response.

DNSSEC - Measurements

“Exploring the overhead of dnssec” [1]:

- simulation based on measurement data and collected traces
- for RSA-signatures **packet size** grows “on average by a factor of 3.4 and 12.7 in the worst case.”
- can lead to a fallback to TCP for the query, adding 5 rtt
- **memory requirements** of nameservers are significantly higher, factor of four for caching NS

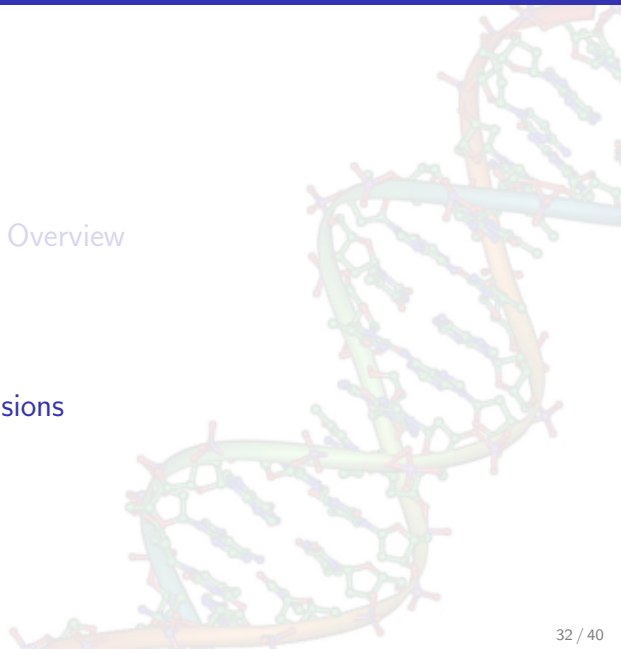
Also:

- home-routers: overflow of resources was observed³
- impact on e.g. root servers is not yet totally clear

³<https://lists.dns-oarc.net/pipermail/dns-operations/2010-September/006123.html>

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Summary - Caching

Effects of caching had been incorrectly estimated:

- main effect is not caching of positive A-answers (heavy tail!)
- help to avoid unnecessary lookups to find authoritative servers
- help by using negative caching to prevent repeated queries
- small TTLs as used by CDNs have no big impact
- positive side-effect of caching: robustness against attacks and misconfigured systems.

Summary - Stability

- DNS is more stable than expected.
- large amount of bogus traffic is not affecting the DNS as a whole due to working delegation and caching
- But: single incidents like the outage of RIPvEs reverse DNS or Microsofts nameservers show, how a configuration error on one site still can affect the system as a whole.

Summary - Scalability

- more and more services depend on DNS and use it as a general purpose data-store or lookup service
- DNS as global address mapping service is working stable and still scaling well, partitioning into independently administered zones works

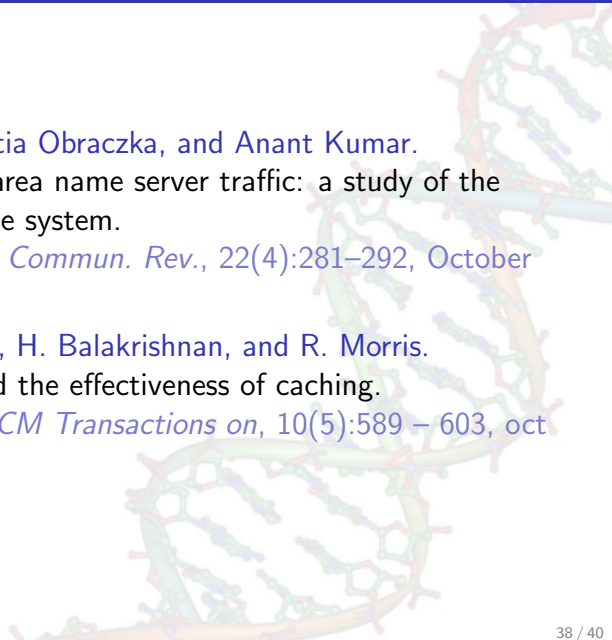


Summary - Issues

- amount of bogus traffic on the rootservers is still surprisingly high and should be continued to act upon.
- The exact effect of extensions like DNSSEC on the load of caching servers and the rootservers is unknown.
- worst case: overloading the servers, denial of service
- impact of initiatives like “A Faster Internet” can currently not be estimated, as there is no data publicly available.


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Download:

<http://f-streibelt.de/talks/>

Acknowledgments:

C.Fuerst for his reliable (email) reminders...

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Image sources:

DNA: [http://en.wikipedia.org/wiki/File:DNA_](http://en.wikipedia.org/wiki/File:DNA_Overview2.png)

[Overview2.png](http://en.wikipedia.org/wiki/File:DNA_Overview2.png)

others: own work and <http://openclipart.org/>

:: ADDITIONAL SECTION:

e.root-servers.net.	85189	IN	A	192.30.140.2
f.root-servers.net.	85189	IN	A	192.5.5.241
c.root-servers.net.	85189	IN	A	192.33.4.12
d.root-servers.net.	85189	IN	A	128.8.10.90
j.root-servers.net.	85189	IN	A	192.58.128.30
k.root-servers.net.	85189	IN	A	193.8.14.129
l.root-servers.net.	85189	IN	A	199.7.83.42
a.root-servers.net.	84603	IN	A	198.41.0.4
b.root-servers.net.	85189	IN	A	192.228.79.281
g.root-servers.net.	85189	IN	A	192.112.36.4
h.root-servers.net.	85189	IN	A	128.63.2.53

```
tcpdump -n -i eth0 -s 0 port 5
tcpdump: listening on eth0, link
00:00:00:00:00:00 IP (tos 0x0, tt
172.16.1.100.43779 > 8.8.8.
00:25:26.734746 IP (tos 0x0, tt
8.8.8.8.53 > 172.16.1.100.43
```